MINIX James

Why Micro-kernel

- Kernel bugs are usually fatal
- User-mode bugs are not

History

- MINIX 1.0 (1987)
 - FOR IBM PC/AT microcomputers
- MINIX 2.0 (1997)
 - Available for x86 and SPARC architecture
 - Follow POSIX
 - Support 32-bit mode
 - TCP/IP stack
 - In kernel drivers
- MINIX 3.0 (2005)
 - Latest version
 - High reliability
 - Many software: X11, emacs, gcc, ftp, etc.

MINIX 3

- Only has a tiny and trusted micro-kernel
- Constantly monitor the status of OS
- Automatically replace or shut down malfunctioning component

Self-repairing OS

- Major faults are properly isolated
- Defects are detected
- Failing components can be replaced on the fly

Architecture



What're in micro-kernel

- Interrupt handling
- Memory management unit (MMU)
- Scheduling
- Interprocess communication (IPC)
- Kernel calls
- Clock driver
- System tasks

What're in user mode

- Drivers
 - Send messages to other processes
 - Make kernel calls (read/write IO ports)
- Servers
 - Network service
 - Has complete TCP/IP stack
 - BSD socket
 - File system service
 - Process user requests
 - Talk to the disk driver for data
 - Call system calls to move data from FS server to user
 - Process management
 - Process creation
 - Memory allocation
 - Handle signals
 - Implement policies
 - Call kernel calls for low-level processing

More servers

- Reincarnation
 - Parent process of all drivers and servers
 - Monitor system well-being by periodically inquires
 - Assumptions
 - No malicious processes
 - Problems can be detected
 - Restarting component could repair a defect
- Data store
 - Tiny database server
 - Backup states
 - Publish system configuration information



MINIX3 features

- 1. Small kernel size
 - 4k lines
- 2. Reduce the damage a bug can do
- 3. Drivers can no longer use privileged calls
 - e.g. disabling interrupt
- 4. Self-repairing
- 5. Soft interrupt and messages
- 6. Need MMU to protect private address space
- 7. Instruction and data space are split
- 8. Scheduling with dynamic priority